

## The efficient congruency method

Again, we consider the system

$$\sum_{i=1}^s x_i^j = \sum_{i=1}^s y_i^j \quad \text{for } 1 \leq j \leq k \quad (5)$$

then given  $\lambda \neq 0, u$ ,  $(x_i, y_i)$  is a solution to (5) if and only if they are

solutions to

$$\sum_{i=1}^s (\lambda x_i + u)^j = \sum_{i=1}^s (\lambda y_i + u)^j \quad \text{for } 1 \leq j \leq k \quad (6)$$

Note this means that  $\underline{z} = (x, y)$  is a solution

$$\Leftrightarrow \lambda \underline{z} + u = (\lambda x + u, \lambda y + u)$$

is a solution

so the system is translation-dilation invariant

Indeed, trivially  $(x, y)$  is a solution  $\Leftrightarrow (\lambda x, \lambda y)$  is a solution

Moreover

$$\sum_{i=1}^s ((x_i + u)^j - (y_i + u)^j) = \sum_{l=1}^{j-1} \binom{j}{l} u^{j-l} \sum_{i=1}^s (x_i^l - y_i^l)$$

(and vice-versa  $x_i = \lambda x_i + u - u$ )  $\Rightarrow (x + u, y + u)$  is a solution

This implies the crucial observation: the number of

solutions  $x_1, \dots, x_s$  to  $\sum_{i=1}^s x_i^j = \sum_{i=1}^s x_{i+s}^j$ , for  $1 \leq j \leq k$   
subject to  $x_i \equiv a \pmod{q} \quad \forall 1 \leq i \leq 2s$

is the number of  $z_1, z_1, \dots, z_s$ ;  $z_i \in X/q$

s.t.  $\sum_{i=1}^s z_i^j = \sum_{i=1}^s z_{i+s}^j$ , for  $1 \leq j \leq k$

which is  $J_{k,s}(X/q)$ .

Initial set-up and notation

$$S(a) = \sum_{1 \leq x \leq X} e(\alpha_1 x + \alpha_2 x^2 + \dots + \alpha_k x^k)$$

Given  $a, \beta \in \mathbb{N}$ ,  $p \in \mathbb{P}$ , write

$$S_a(a, \beta) = \sum_{\substack{1 \leq x \leq X \\ x \equiv \beta \pmod{p^a}}} e(\alpha_1 x + \dots + \alpha_k x^k) \quad (7)$$

We'll fix a prime  $p > k$ , but will vary  $\beta$  and  $a$ .

We'll count solutions to the Vinogradov system

$$x_1^j + \dots + x_s^j = x_{s+1}^j + \dots + x_{2s}^j, \quad 1 \leq j \leq k \quad (8)$$

subject to the congruence conditions, namely

ANDE XIII (3)

$$X_i \equiv \xi \pmod{p^a} \text{ for } 1 \leq i \leq m \text{ and } s+1 \leq i \leq s+m, \quad (8)$$

$$X_i \equiv \eta \pmod{p^b} \text{ for } m+1 \leq i \leq s \text{ and } s+m+1 \leq i \leq 2s$$

which we'll denote by  $I_m(X; \xi, \eta, a, b)$ .

Note that

$$I_m(X; \xi, \eta, a, b) = \int_{[0,1]^k} |S_a(\underline{x}, \xi)|^{2m} |S_b(\underline{x}, \eta)|^{2(s-m)} d\underline{x} \quad (10)$$

Indeed, the RHS of (10) is

$$\begin{aligned} & \int_{[0,1]^k} \left| \sum_{\substack{\chi \leq X \\ \chi \equiv \xi \pmod{p^a}}} e\left(\sum_{j=1}^k \alpha_j \chi^j\right) \right|^{2m} \left| \sum_{\substack{\chi \leq X \\ \chi \equiv \eta \pmod{p^b}}} e\left(\sum_{j=1}^k \alpha_j \chi^j\right) \right|^{2(s-m)} d\underline{x} \\ &= \int_{[0,1]^k} \sum_{\substack{\chi_1, \dots, \chi_m \\ \chi_{s+1}, \dots, \chi_{s+m} \\ \equiv \xi \pmod{p^a}}} \sum_{\substack{\chi_{m+1}, \dots, \chi_s \\ \chi_{m+s+1}, \dots, \chi_{2s} \\ \equiv \eta \pmod{p^b}}} e\left(\sum_{j=1}^k \alpha_j (\chi_1^j + \dots + \chi_m^j + \chi_{m+1}^j + \dots + \chi_s^j - \chi_{m+s+1}^j - \dots - \chi_{2s}^j)\right) \end{aligned}$$

which is the number of solutions to (8)-(9).

If  $m=0$  then the expression in (10) is independent of  $\xi \in a$ ; and we write

$$I_0(X; \eta, b) = \int_{[0,1]^k} |S_b(\underline{a}, \eta)|^{2s} d\underline{a} \quad (11)$$

and  $I_0(X; b) = \max_{\eta \pmod{p}} I_0(X; \eta, b)$

similarly for  $1 \leq m \leq s-1$ , write

$$I_m(X; a, b) = \max_{\eta \not\equiv \xi \pmod{p}} I_m(X; \xi, \eta, a, b)$$

The classical p-adic method (Linnik, Kacivabian)

Thm (Vinogradov) Let  $s > s_k = \frac{1}{2}k(k+1)$ . Then one has

$$J_{k,s}(X) \ll_{k,s} X^{2s - s_k + \eta_{s,k}}; \text{ with } \eta_{s,k} \leq \frac{k^2}{2} \left(1 - \frac{1}{k}\right)^{\lfloor s/k \rfloor} \quad (12)$$

Note • If  $s \geq ck^2 \log k$  then  $\eta_{s,k} \leq \frac{k^2}{2} \left(1 - \frac{1}{k}\right)^{-ck \log k} \leq k^{-c}$

and in this range the estimate is almost as good as the main conjecture.

• We'll only sketch the "p-adic" proof of Linnik and Karatsuba.

Lemma. Let  $p \leq X$ ,  $l \geq 1$ . Then  $J_{k, k+l}(X) \leq p^{2l} I_k(X; p, 1)$  (13)

Pf Recall that  $I_k(X; 1) = \sup_{\mathfrak{z} \pmod{p}} \int_{[0, 1]^k} |S_1(\underline{a}, \mathfrak{z})|^{2l} |S(\underline{a})|^{2k} d\underline{a}$

Also,  $|S(\underline{a})|^{2l} = \left| \sum_{\mathfrak{z} \pmod{p}} S_1(\underline{a}, \mathfrak{z}) \right|^{2l} \leq p^{2l-1} \sum_{\mathfrak{z} \pmod{p}} |S_1(\underline{a}, \mathfrak{z})|^{2l}$

$$\begin{aligned} \Rightarrow J_{k, k+l}(X) &\leq p^{2l-1} \sum_{\mathfrak{z} \pmod{p}} \int_{[0, 1]^k} |S(\underline{a})|^{2k} |S_1(\underline{a}, \mathfrak{z})|^{2l} d\underline{a} \\ &\leq p^{2l} \sup_{\mathfrak{z} \pmod{p}} \int_{[0, 1]^k} |S(\underline{a})|^{2k} |S_1(\underline{a}, \mathfrak{z})|^{2l} d\underline{a} \end{aligned}$$

Note We have only used Hölder's ineq. □

i.e.  $\frac{1}{p} S(\underline{a}) = \frac{1}{p} \sum_{1 \leq x \leq X} e(\alpha_1 x + \dots + \alpha_k x^k) = \frac{1}{p} \sum_{\mathfrak{z} \pmod{p}} \sum_{\substack{1 \leq x \leq X \\ X \equiv \mathfrak{z} \pmod{p}}} e(\alpha_1 x + \dots + \alpha_k x^k)$

$$\Rightarrow \left( \frac{1}{p} |S(\underline{a})| \right)^{2l} \leq \frac{1}{p} \sum_{\mathfrak{z} \pmod{p}} |S_1(\underline{a}, \mathfrak{z})|^{2l} = \frac{1}{p} \sum_{\mathfrak{z}} S_1(\underline{a}, \mathfrak{z})^{2l}$$

- We've embedded a congruence relation into the count  $J_{s,k+l}(X)$ .

We use translation/dilation invariance to compare

$I_k(X; 0, i)$  to  $J_{s,k}(X/p)$  to obtain a recursive inequality.

Lemma 2. (Linnik's lemma) Let  $p \in \mathbb{P}$  s.t.  $p^2 > X$ ,  $n_1, \dots, n_k \in \mathbb{Z}$ .

Then  $(X_1, \dots, X_k)$  such that  $1 \leq X_i \leq p^k$  satisfying

$$S_j(x) = \sum_{i=1}^k X_i^{j_i} \equiv n_j \pmod{p^j}, \text{ for } \forall 1 \leq j \leq k$$

is at most  $k! p^{\frac{1}{2}k(k-1)}$

Pf: Given  $1 \leq j \leq k$  there are  $p^{k-j}$  values of  $1 \leq w_j \leq p^k$

s.t.  $w_j \equiv n_j \pmod{p^j}$ . Given  $1 \leq w_1, \dots, w_k \leq p^k$

$$\text{st } \sum_{i=1}^k X_i^{j_i} \equiv m_j \pmod{p^j} \quad \text{and} \quad \left. \begin{array}{l} 1 \leq j \leq k \\ 1 \leq y_i \leq p^k \end{array} \right\}$$

$$\sum_{i=1}^k y_i^{j_i} \equiv w_j \pmod{p^j}$$

$$\text{Hence } \prod_{i=1}^k (z - X_i) \equiv P(z) \equiv \prod_{i=1}^k (z - y_i) \Rightarrow P(y_i) \equiv 0 \pmod{p^k}$$

$\Rightarrow y_1, \dots, y_k$  is a permutation of  $X_1, \dots, X_k$ .  $\square$

Lemma 3. We have  $I_{k,2+l}(X;0,1) \leq k! p^{\frac{k+l-1}{2}} X^k J_{k,l}(X/p)$  (7)

Pf By definition  $I_{k,2+l}(X;0,1)$  is the max of the number solutions  $(x_1, \dots, x_k, y_1, \dots, y_l, u_1, \dots, u_l, v_1, \dots, v_l) \in \mathbb{Z}$

$$(15) \quad \sum_{i=1}^k (x_i^j - y_i^j) = \sum_{m=1}^l ((p u_m + 3)^j - (p v_m + 3)^j) \quad ; 1 \leq j \leq l$$

as it is the max. (over  $\mathbb{Z} \pmod{p}$ ) of the number of solutions to the Vinogradov system

$$(16) \quad \sum_{i=1}^k x_i^j + \sum_{m=1}^l u_m^j = \sum_{i=1}^k y_i^j + \sum_{m=1}^l v_m^j \quad ; 1 \leq j \leq l$$

under the congruence restrictions  $u_m \equiv v_m \equiv 3 \pmod{p}$  for  $1 \leq m \leq l$ .

By translation invariance (15) is equivalent to:

$$\sum_{i=1}^k ((x_i - 3)^j - (y_i - 3)^j) = p^j \sum_{m=1}^l (u_m^j - v_m^j) \quad (17)$$

We solve this in 2 steps, first count the no. of sol's

$$\text{to:} \quad \sum_{i=1}^k (x_i - 3)^j \equiv \sum_{i=1}^k (y_i - 3)^j \pmod{p^j} \quad (18)$$

(2)

One may choose  $x_{1-1} y_k$  via  $X^k$  drags; and once  $x_{1-1} y_k$  fixed, by Linnik's Lemma using  $X \leq p^k$  there are at most  $e! p^{\frac{1}{2}k(k-1)}$  choices for  $y_{1-1} y_k$ .

Writing  $p^i w_j = \sum_{i=1}^k ((x_i - \frac{1}{3})^k - (y_i - \frac{1}{3})^k)$ , there

are at most  $J_{k,e}(X/p)$  choices for  $1 \leq u_m, v_m \leq X/p$

solving 
$$\sum_{m=1}^l u_m^j - v_m^j = w_j, \quad 1 \leq j \leq k$$

uniformly in  $w_j$ . Indeed, the # of solutions

may be written as

$$\int_{[0,1]^k} \left| \sum_{\substack{u \in \mathbb{Z} \\ |u| \leq X/p}} e(\alpha_1 u + \alpha_2 u^2 + \dots + \alpha_k u^k) \right|^{2l} e(-\alpha_1 w_1 - \dots - \alpha_k w_k) d\alpha$$

$$\leq \int_{[0,1]^k} |S_k(\alpha, X/p)|^{2l} = J_{k,e}(X/p)$$

□